

Desarrollo de Aplicaciones en Red

José Rafael Rojano Cáceres
<http://www.uv.mx/rrojano>

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What we saw



- **Distributed system:** Collection of independent computers that for the user works like if it were one CPU.
- **Distributed operating system:** Is the actor in charge to achieve the past.
- **Network operating systems:** is the actual result, wait to to became in a true DOS. Each machine have their kernel.

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1. OS Design

- The design of OS can be saw from three points of view:
 - *What services are offered*
 - *Interface available for users and programmers*
 - *Component and interconnections*

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1.1 Operating System Service (1)

- The number and type of service differs between each OS. The basic interface offered are:
 - **User interface**
 - *CLI, GUI, and batch interface*
 - **Program execution**
 - **I/O operations**
 - *ports or device*
 - **File-system manipulation**
 - **Communications**
 - *Shared memory or message passing*
 - **Error detection**

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1.1 Operating System Service (2)

- **Secondary services can be found**
 - **Resource allocation**
 - *CPU scheduling routines according to the complexity of the job*
 - **Accounting**
 - **Protection and security**

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2. Service description

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2.1 User interface (1)

- As we saw, CLI a GUI approach exists
- CLI are operated by a **command interpreter**
- The command interpreter is called **shell**
- Two approach for building a interpreter
 - *by itself execute programs (size of interpreter)*
 - *Implements command as a system programs (to load and execute programs)*

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2.1 User interface (2)

- GUI started in the early of 70s at Xerox
- Widespread by the Macintosh in the 80s
- Several GUIs does exist
 - *Mac OS*
 - *Windows*
 - *CDE and X-window*
 - *KDE, GNome*

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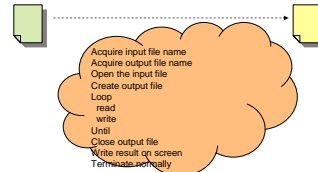
3. System calls



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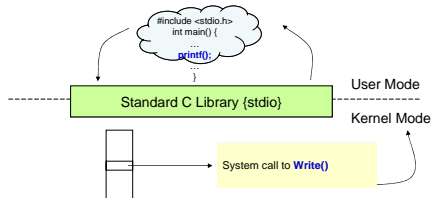
3.1 System calls

- It provides an interface for calling services in an OS
- *Imagine how many system calls does exist in order to copy from one file to another*
- System programmers usually consider this aspect to provide an **API** (*application programming interface*) for the developer



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- For APIs use does exist different ones for example:
 - *Win32*
 - *Java*
 - *Posix API for posix systems includes Linux, Mac, Unix*
- Also it's important to say that many times the programmer use API to do system calls in stead of invoke system calls, this help to the **portability**



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3.2 Type of System Calls

- Five general group are found:
 - *Process control*
 - *File manipulation*
 - *Device manipulation*
 - *Information maintenance*
 - *Communications*

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3.2.1 Process control

- Process control refers to the mechanism that OS use for handling the termination of a program.
- It can be normal (end) or abnormal (abort).
- In such case of error, a trap is issue causing dump file.
- In the next slide you will see a resume of type of system call

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- Process control
 - End, abort
 - Load, execute
 - Create or terminate process
 - Get or set process attribute
 - Wait
 - Allocate memory
- File management
 - Create, delete, Open, close, read, write
- Device management
 - Request, release a device. Read or write
- Information maintenance
 - Get or set, time, system data or process properties
- Communications
 - Create, delete connection, send or receive message

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Example of getting process properties

```

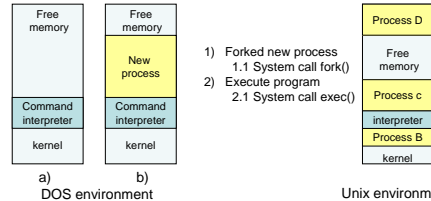
1.  #include <sys/types.h>
2.  #include <stdio.h>
3.  #include <unistd.h>
4.  void main(void)
5.  {
6.      pid_t id_proceso;
7.      pid_t id_padre;
8.      id_proceso = getpid();
9.      id_padre = getppid();
10.     printf("Identificador de proceso: %d\n", id_proceso);
11.     printf("Identificador del proceso padre: %d\n",
12.         id_padre);

```

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3.2.1.1 Multitasking environment

- In multitasking environment in stead of single-tasking child process can be executed. Mechanism are different, Ex.:



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Example of create process

```

1.  #include <sys/types.h>
2.  #include <stdio.h>
3.  #include <unistd.h>
4.  void main(void)
5.  {
6.      pid_t pid;
7.      int status;
8.      pid = fork();
9.      switch(pid)
10.     {
11.         case -1: /* error del fork() */
12.             perror("Error al el proceso fork");
13.             break;
14.         case 0: /* proceso hijo */
15.             execlp("ls", "ls", "-l", NULL);
16.             perror("Error en el proceso de creación exec");
17.             break;
18.         default: /* padre */
19.             printf("Yo soy el proceso padre\n");
20.     }
21. }

```

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3.2.2 File management

- The design of File System has many aspect, however we can identify some common system calls in this aspect:
 - Basic operations over files (O, C, R, W, S)
 - Get or set attributes to files and directories

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3.2.3 Device management

- One a process is executed in need access to different resource, the resource controlled by the system can be seen as device
- Device can be hardware or software
 - *Disk, printers*
 - *Virtual devices*
- So, in this category is necessary to consider the administration for accessing and releasing devices

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3.2.4 Information maintenance

- In this aspect many systems provide functionality for date and time.
- Also the system provides mechanism to identify process, as we saw in the example on page 15.

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3.2.5 Communication

- As we said there are two ways for
 - *Shared memory*
 - *Message passing*

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3.2.5.1 Shared memory

- *To use this model the process employ system calls to access shared memory owned by other process*
- *Shared memory requires that process removes or grant access level*
- *It is necessary for this model that programs provide mechanism for concurrency*
- *There is a variant to this model, thread, where shared memory is the default mechanism*

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3.2.5.2 Message passing

- In this model through message are send between process
- Communication must be establish in advanced
- Name of the other process and communicator (PC) must be known
- Identify this properties through system calls are sent to the correspondent system calls for processing

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3.3 System Programs

- *It's another aspect to consider in the design of an OS*
- *Many system programs are just interface of the system calls*
- *They are divided in the next categories*
 - File management
 - Status information
 - File modification
 - Programming language support
 - Programming loading and execution
 - communications

Programación de sistemas

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4. Design & implementation in OS

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General aspects

- Define goals and specification at highest level: considering hardware
- But, beyond this consideration basically we have two approach
 - User goals
 - System goals
- In general we found the rules of design in **software engineering**

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5. Operating System Structure

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5.1 General aspects

- The approach to design the structure of an OS is in general based in the **partition task** this is because of complexity

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5.2 Simple structure

- Many commercial OS does not have defined structures, because are small like Msdos
- In Msdos interfaces and levels are not well separated for that applications can access basic functionalities of I/O
- Another example is the early Unix which was separated in kernel and system programs, later the kernel was layered in services (see fig 5.2)

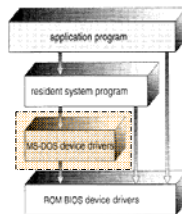


Fig. 5.1 MS-DOS layer structure.

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5.3 Layered structure (1)

- With proper hardware support, OS can be broken into smaller and efficient pieces.
- In the figure 5.2 and 5.3 you can see an approach where OS is broken in layers or levels. The bottom (0) is hardware and upper level (N) are applications
- The main advantage in this approach is simplicity and debugging
- Each layer is implemented with those functions provides for the lower layer
- A layer does not need to know how are operations implemented, only need to know what it does

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5.3 Layered structure (2)

- Also layered approach can be less efficient than others types
- For example a system call can take so much time going from upper to lower considering that it may need execute I/O operations, the system call can be longer than in a non layered system.

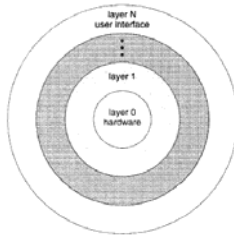


Fig. 5.2 A layered operating system.

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Unix systems structure

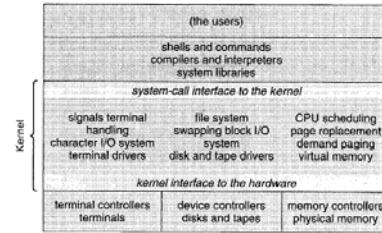


Fig. 5.3 UNIX system structure.

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5.4 Microkernel

- As was said the expansion of kernel in UNIX become complex, in 80s, Carnegie Mellon University develop *Mach system* with an approach in microkernel.
- This method consist in **remove all nonessential components** from the kernel implementing them as system or user-level programs.
- The main function is to provide an interface between user and services through **message passing**.
- Modularity and easily modifications without affect the kernel
- There is also more security because programs are executed with user privileges
- Ex.: Tru64 unix, QNX

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5.5 Modules

- Another approach is the construction of OS as modules
- This approach is based in the object oriented paradigm
- Here the kernel has some core components and dynamically links to services
- Examples of this are Solaris, Linux and MacOS

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A core kernel structure



Fig. 5.4 Solaris loadable modules.

- **Message passing does not exist**

- MacOS knows as Darwin as hybrid: layers and microkernel

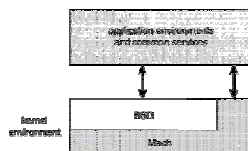


Fig. 5.5 The Mac OS X structure.

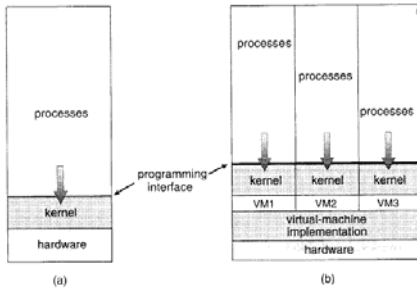
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Virtual Machines

- A virtual machine is the conclusion of a layered system.
- The idea in a VM is the **abstraction** of hardware into several different execution environments, creating the idea of a own computer.
 - Virtual memory
 - Scheduling

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Model of VM



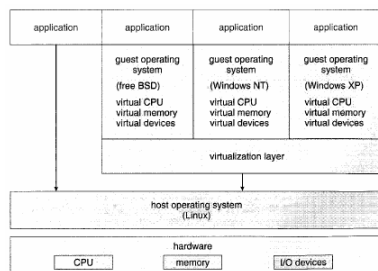
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VMware example

- VMware abstract Intel x86 hardware into isolated VM.
- It let run different guest OS

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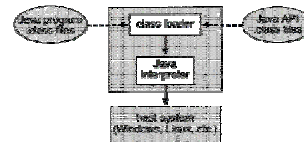
VMware architecture



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Java VM example

- Java defines an architectural-neutral byte code
- Two components for java architecture



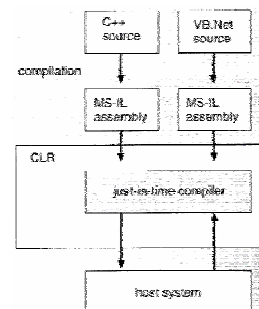
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.Net platform example

- Also program are targeted to run on platform instead of specific architecture
- The core of the architecture is the CLR (Common Language Runtime)
 - Which run intermediate code representation call MS-IL (Microsoft Intermediate Language)

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.Net architecture



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Reference

- Silberschatz et al, Operating Systems concepts 7th. Wiley.

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